

# On Maximum Differential Graph Coloring

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**Abstract.** We study the *maximum differential graph coloring problem*, in which the goal is to find a vertex labeling for a given undirected graph that maximizes the label difference along the edges. This problem has its origin in map coloring, where not all countries are necessarily contiguous. We define the differential chromatic number and establish the equivalence of the maximum differential coloring problem to that of  $k$ -Hamiltonian path. As computing the maximum differential coloring is NP-Complete, we describe an exact backtracking algorithm and a spectral-based heuristic. We also discuss lower bounds and upper bounds for the differential chromatic number for several classes of graphs.

## 1 Introduction

The Four Color Theorem states that only four colors are needed to color any map so that no neighboring countries share the same color. This theorem assumes that each country forms a contiguous region in the map. However, if countries in the map are not all contiguous then the result no longer holds [6]. Instead, this necessitates use of a unique color for each country to avoid ambiguity. As a result, the number of colors needed is equal to the number of countries.

Given a map, we define the country graph  $G = \{V, E\}$  to be the undirected graph where countries are nodes and two countries are connected by an edge if they share a nontrivial boundary. We then consider the problem of assigning colors to nodes of  $G$  so that the color distance between nodes that share an edge is maximized. Figure 1 gives an illustration of this map coloring problem.

As not all colors make suitable choices for country colors, a good color palette is often a gradation of certain “map-like colors”. Furthermore, if the final result is to be printed in black and white, then the color space is strictly 1D. We therefore make the assumption that the colors form a line in the color space, and model our map coloring problem as one of node labeling of the graph, where the available labels are from set  $C = \{\text{all permutations of } \{1, 2, \dots, |V|\}\}$ . This brings us to the *maximum differential graph coloring problem*, in which we aim to find a vertex labeling for a given undirected graph that maximizes the label difference along the edges in the graph.

More formally, we are looking for a labeling function, a bijection  $c : V \rightarrow \{1, 2, \dots, |V|\}$ , that solves the following MaxMin optimization problem:

$$\max_{c \in C} \min_{\{i,j\} \in E} w_{ij} |c(i) - c(j)| \quad (1)$$

Here  $w_{ij}$  are positive weights representing importance of keeping the difference of labels between nodes  $i$  and  $j$  large. To simplify the problem further, throughout this paper we assume  $w_{ij} = 1$ .

Figure 2 illustrates the graph coloring problem. We use a color palette of yellow to blue, indexed from 1 to 9. Figure 2(a) shows a map with 9 countries, with one non-contiguous country (its two components are at the center, and to the far right of the map). Figure 2(b) shows the corresponding country graph. The node labeling is not optimal, with many adjacent nodes having label difference of 1. Figure 2(c) gives the optimal node labeling, with a minimal label difference of 3. The map in Figure 2(a) is in fact colored using this optimal coloring scheme, which gives distinctive colors for neighboring countries.

In this paper we define the differential chromatic number and show a correspondence between the maximum differential graph coloring problem and the Hamiltonian path problem in the complement graph. We also provide exact and heuristic algorithms for computing good solutions. In Section 3 we establish the relationship between this problem and that of finding a  $k$ -Hamiltonian path. a heuristic algorithm on a number of well known graphs. Section 5 gives results for some special graphs. We conclude the paper with Section 6.

## 2 Related Work

The problem of maximum differential coloring of graphs arises in the context of coloring a map in which not all regions are necessarily contiguous [6, 7]. A variation of the differential graph coloring problem was studied by Dillencourt *et al.* [4], under the assumption that all colors in the color spectrum are available. This makes the problem continuous rather than discrete. A heuristic algorithm based on the force-directed model is used to select  $|V|$  colors as far apart as possible in the 3-dimensional color space. However this algorithm cannot be used directly for the general map coloring problem, as maps are typically colored with “map-like colors”, often light, pastel colors which come from a very restricted subset of the color spectrum. It may be possible to adapt the algorithm and apply it to a lower-dimensional color manifold, but because the algorithm is greedy it is more likely to converge to local minimum in lower dimensional space.

The problem of finding a permutation that *minimizes* the labeling differences along the edges is well-studied problem in the context of minimum *bandwidth* or wavefront reduction ordering for sparse matrices. It is known that the bandwidth problem is NP-complete, with a reduction from 3SAT dating to 1975 [11]. Moreover, it is NP-complete to find any constant approximation, even when restricting the problem to trees [1, 15]. A number of effective heuristics for that problem have been proposed [8, 9, 13].

The complement of the bandwidth problem is that of *maximizing* the labeling difference along the edges and is less well known as the *antibandwidth* problem.

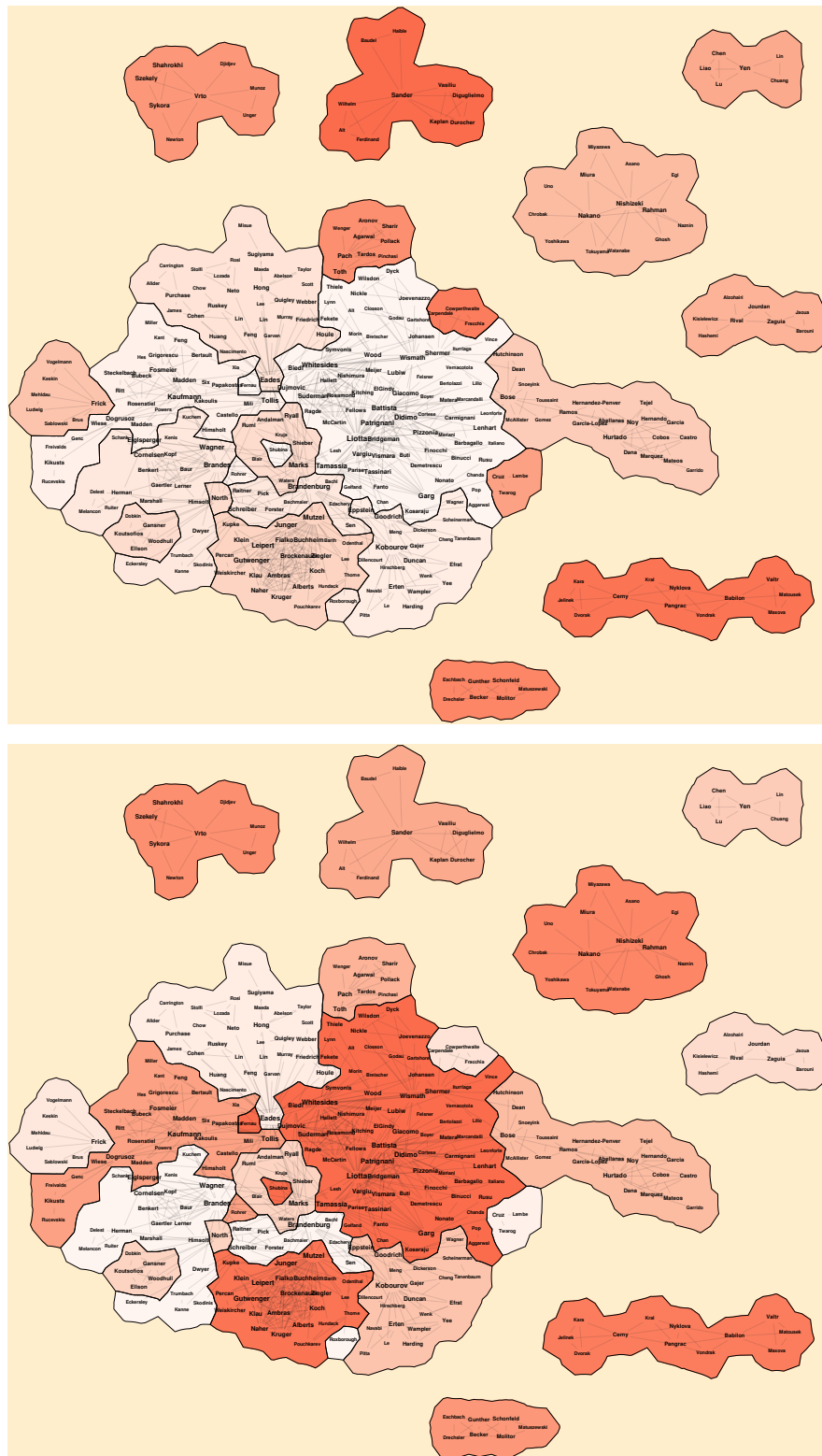


Fig. 1. Graph Drawing Symposia (1994-2004) co-authorship map with a single-hue color palette from ColorBrewer [2]: a random color assignment above and optimal below.

Not surprisingly, this is also an NP-Complete problem as shown in 1984 [10]. Exact values for the antibandwidth of some of Hamming graphs [5], as well as for meshes [14], hypercubes [12, 16], and complete  $k$ -ary trees for odd values of  $k$  [3].

There have been definitions for a graph being  $k$ -edge Hamiltonian and  $k$ -vertex Hamiltonian in the literature [17], which are defined as the graph having a Hamiltonian cycle after removing any  $k$  edges or vertices, respectively. These concepts, though related, are different from our definition of  $k$ -Hamiltonian path, introduced in the next section.

### 3 Equivalence between Maximum Differential Coloring and Hamiltonian Path

Let the undirected graph of interest is  $G = \{V, E\}$ . We denote  $\bar{G}$  as the complement of graph  $G$ , defined as  $\bar{G} = \{V, \bar{E}\}$ , where  $\bar{E} = \{\{i, j\} | i \neq j, i, j \in V, \text{ and } \{i, j\} \notin E\}$ . In other words,  $\bar{G}$  is the graph containing all nodes in  $G$ , and all edges that are not in  $G$ . We now define the maximum differential coloring problem formally.

**Definition 1.** A coloring of the nodes of  $G$  is a bijection  $c : V \rightarrow \{1, 2, \dots, |V|\}$ . We denote the set of all colorings  $C(G)$ .

**Definition 2.** A  $k$ -differential coloring of  $G$  is one in which the absolute coloring difference of the endpoints for any edge is  $k$  or more. We denote the set of all such  $k$ -differential colorings  $DC(G, k) = \{c \mid c \in C(G), |c(i) - c(j)| \geq k \text{ for all } \{i, j\} \in E\}$ .

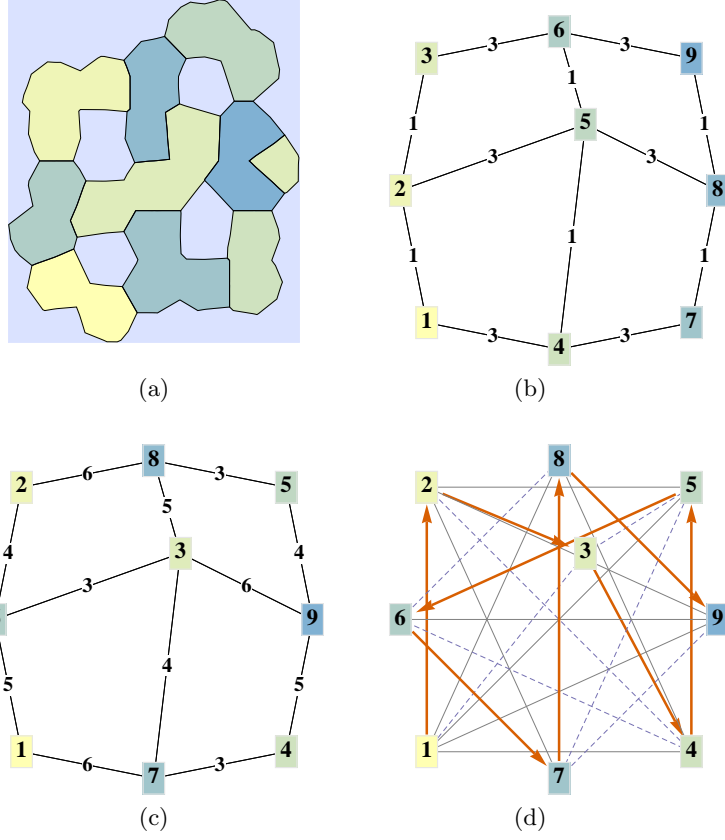
**Definition 3.** A graph is  $k$ -differential colorable if  $DC(G, k) \neq \emptyset$ .

**Definition 4.** If a graph is  $k$ -differential colorable, but not  $(k + 1)$ -differential colorable, it is said to have a differential chromatic number  $k$ , and denoted as  $dc(G) = k$ .

**Definition 5.** A Hamiltonian path of  $G$  is a bijection  $p : \{1, 2, \dots, |V|\} \rightarrow V$ , such that  $\{p(i), p(i + 1)\} \in E$  for all  $i = 1, 2, \dots, |V| - 1$ . We denote the set of all Hamiltonian paths  $H(G)$ .

The key insight in understanding the maximum differential coloring comes from observing a good differential coloring scheme. For example, in Figure 2(c), the minimum labeling difference between any two adjacency nodes is 3. This means that any two nodes with labeling difference of  $\leq 2$  can not form an edge in  $G$ , in other word they must form an edge in the complement of the graph, shown in Figure 2(d). Therefore the list of nodes induced by this coloring scheme of  $G$ , with labels  $\{1, 2, \dots, |V|\}$ , forms a Hamiltonian path in  $\bar{G}$ , shown in Figure 2(d) with orange arrows. Furthermore, this Hamiltonian path is such that any two nodes along the path with labeling difference of 2 are also connected by an edge of  $\bar{G}$ , shown in Figure 2(d) with dashed lines.

This observation leads to a natural extension of the concept of Hamiltonian path, which we call  $k$ -Hamiltonian path.



**Fig. 2.** An illustration of the graph coloring problem and its relation to a  $k$ -Hamiltonian path. A yellow to blue coloring scheme (indexed from 1 to 9) is used; nodes are labeled by color indices and edges are labeled with the absolute difference of the adjacent nodes. (a) An input map. (b) The country graph  $G$  corresponding to the input map; note that many adjacent nodes have very similar colors. (c) An optimal 3-differential coloring of the country graph. (d) The complement graph  $\bar{G}$  corresponding to  $G$  and a 2-Hamiltonian path. The Hamiltonian path  $\{1, 2, \dots, 9\}$  is shown with orange arrows, and edges with labeling difference of 2 are shown as blue dashed lines. Note that the 2-Hamiltonian path of  $\bar{G}$  is such that any two nodes of labeling difference  $\leq 2$  form an edge of  $\bar{G}$ .

**Definition 6.** A  $k$ -Hamiltonian path ( $k \geq 1$ ) of  $G$  is a Hamiltonian path, such that each  $i$ -th node on the path is connected to the  $j$ -th node, if  $|i - j| \leq k$ . We define the set of  $k$ -Hamiltonian paths as  $H(G, k) = \{p | p \in H(G), \text{ and } \{p(i), p(j)\} \in E \text{ if } 0 \leq i, j \leq |V| \text{ and } |i - j| \leq k\}$ . Clearly  $H(G) = H(G, 1)$ .

Based on our previous discussion, we can relate the  $k$ -differential coloring problem to that of finding a  $(k - 1)$ -Hamiltonian path.

**Theorem 1.** A  $k$ -differential coloring ( $k \geq 2$ ) of a graph  $G$  exists if and only if a  $(k - 1)$ -Hamiltonian path exists in the complement of  $G$ , i.e.,  $DC(G, k) \neq \emptyset$  if and only if  $H(\bar{G}, k - 1) \neq \emptyset$ . Furthermore, the inverse function of each element of  $DC(G, k)$  is in  $H(\bar{G}, k - 1)$ , and vice visa.

*Proof.* Suppose a  $k$ -differential coloring  $c$  exists for  $G$ . Define the path,  $p = c^{-1}$ , so that it visits the vertices in the graph in order of their color index (i.e., the first vertex is that with color 1, the second is that with color 2 and so on). Consider two nodes  $u = p(i)$  and  $v = p(j)$  such that  $|i - j| \leq k - 1$ . Since the color difference of these two nodes in the original graph,  $|c(u) - c(v)| = |i - j| \leq k - 1$ , by definition of  $k$ -differential coloring  $(u, v)$  is not an edge of  $G$ , hence  $\{u, v\} = \{p(i), p(j)\}$  is an edge of  $\bar{G}$ . It follows that  $p = c^{-1}$  is a  $(k - 1)$ -Hamiltonian path.

Conversely, suppose a  $(k - 1)$ -Hamiltonian path  $p$  exists for  $\bar{G}$ . Define a coloring  $c = p^{-1}$ , where the color index is given by the order in which a vertex appears in the path (i.e., color 1 is assigned to the first vertex along the path, color 2 to the second, and so on). Consider any edge  $\{u, v\} \in E$ . We prove that  $|c(u) - c(v)| \geq k$ . Assume that  $|c(u) - c(v)| < k$ . Let  $i = c(u)$  and  $j = c(v)$ , then  $|i - j| < k$ , and  $u = c^{-1}(i) = p(i)$  and  $v = c^{-1}(j) = p(j)$ . By definition of a  $(k - 1)$ -Hamiltonian path,  $(u, v)$  must be an edge of  $\bar{G}$ , which is a contradiction with the fact that  $\{u, v\} \in E$ . It follows that  $|c(u) - c(v)| \geq k$ , and so  $c = p^{-1}$  is a  $k$ -differential coloring of  $G$ .  $\square$

This theorem immediately gives an upper bound for the differential chromatic number of a graph based on its maximum degree:

**Corollary 1.** *A graph of maximum degree  $\Delta(G)$  has a differential chromatic number of at most  $|V| - \Delta(G)$ .*

*Proof.*  $\bar{G}$  must have a node with degree  $|V| - 1 - \Delta(G)$ , therefore  $H(\bar{G}, |V| - \Delta(G)) = \emptyset$ , or by Theorem 1  $DC(G, |V| - \Delta(G) + 1) = \emptyset$ .  $\square$

A special case of differential coloring is finding a scheme with maximum difference of 2 or more. Clearly:

**Corollary 2.** *Finding a 2-differential coloring of a graph  $G$  is equivalent to finding a Hamiltonian path of  $\bar{G}$ .*

Given the equivalence between 2-differential coloring and Hamiltonian path, a number of well known results for Hamiltonian path can immediately be used to give results on 2-differential coloring.

**Theorem 2.** *(Ore's Theorem) A graph with more than 2 nodes is Hamiltonian if, for each pair of non-adjacent nodes, the sum of their degrees is  $|V|$  or greater*

**Corollary 3.** *A graph with more than 2 nodes is 2-differential colorable if, for each pair of adjacent nodes, the sum of their degrees is  $|V| - 1$  or less.*

**Theorem 3.** *(Dirac's Theorem). A graph with more than 2 nodes is Hamiltonian if each node has degree  $|V|/2$  or greater.*

**Corollary 4.** *A graph with more than 2 nodes is 2-differential colorable if each node has degree  $|V|/2 - 1$  or less.*

Corollaries 3-4 confirm the intuition that a sparser graph has a better chance of being more differential colorable.<sup>3</sup> The flip side of this intuition is that the complement graph needs to be denser. For a graph to be  $k$ -colorable, the complement graph must be well connected. The next theorem follows from a result in [12]:

**Theorem 4.** *The complement of a  $k$ -differential colorable graph is  $(k - 1)$ -connected.*

Given the connection between finding a 2-differential coloring and a Hamiltonian path, it is not difficult to prove directly that the maximum differential coloring problem is NP-Complete. An equivalent result was established in 1984 in the context of the antibandwidth problem [10].

## 4 Algorithms for Maximum Differential Coloring

Theorem 1 provides a way to check whether a graph is  $k$ -differential colorable. The following `kpath` algorithm attempts to find a  $k$ -Hamiltonian path of  $\bar{G} = \{V, \bar{E}\}$  from a given node  $i$ . Before calling `kpath` the  $k$ -Hamiltonian path is initialized as  $p = \emptyset$ . Each call to `kpath` recursively tries to add a neighbor to the last node in the path, and checks that this maintains the  $k$ -path condition. If the condition is violated, the next neighbor is explored, or the algorithm backtracks. If a  $k$ -Hamiltonian path of  $\bar{G}$  is found by the algorithm, we know that the graph  $G$  is  $(k + 1)$ -colorable.

This exact algorithm requires exponential time, making it impractical for large graphs, where heuristic algorithms might be a better choice. Gansner *et al.* [6] proposed a heuristic based on a relaxation of the discrete MaxMin problem (1) into a continuous maximization problem of 2-norm:

$$\max \sum_{\{i,j\} \in E} w_{ij}(c_i - c_j)^2, \text{ subject to } \sum_{i \in V} c_i^2 = 1 \quad (2)$$

where  $c \in R^{|V|}$ . This continuous problem is solved when  $c$  is the eigenvector corresponding to the largest eigenvalue of the weighted Laplacian of the graph. Once (2) is solved, the ordering of the eigenvector is used as an approximate solution for the MaxMin problem. This is followed by a greedy refinement algorithm which repeatedly swaps pairs of vertices, provided that the swap improves the coloring scheme. We call this algorithm `GSpectral` (Greedy Spectral). Table 1 gives the chromatic number of some well known graphs found by `kpath`, as well as an estimate of the differential chromatic number given by the `GSpectral` algorithm. The `GSpectral` algorithm gets close to the optimal, though it finds the actual optimal only twice. The exact algorithm `kpath` can be prohibitively

<sup>3</sup> Even planar graphs, which have average degree less than 6, may be tough to color well. Consider a star graph, where one vertex is connected to every other vertex in the graph. Regardless of the coloring, at least one pair of nodes will have a difference of at most one, thus showing that  $dc(G) = 1$ .

**Input:**  $\bar{G}, i, p$  (initialized to  $\emptyset$  on first entry)  
**Output:**  $p$   
**if**  $i \in p$ , *or*  $i$  is not connected with the last  $k$  nodes in  $p$  **then**  
    return  $\emptyset$ ;  
**end**  
 $p = p \cup \{i\}$ ;  
**if**  $|p| = |V|$  **then**  
    return  $p$ ;  
**end**  
**foreach** neighbor node  $j$  of  $i$  **do**  
    **if**  $\text{kpath}(\bar{G}, j, p) \neq \emptyset$  **then**  
        return  $p$ ;  
    **end**  
**end**  
 $p = p - \{i\}$ ;  
return  $\emptyset$ ;

**Algorithm 1:** `kpath`

expensive even for small graphs. For example, on the 60-node football graph (skeleton graph of a truncated icosahedron), one week of CPU time was not enough to find the exact differential chromatic number using `kpath`, whereas the greedy spectral algorithm gives a lower bound of 18 in a few milli-seconds. We also tested `GSpectral` on larger grid graphs for which we believe that we know the differential chromatic number (see Section 5). For `grid10` and `grid20` graph, it gives an estimate of 30 and 124, vs the actual differential chromatic numbers of 45 and 190, respectively. In both case the CPU time for `GSpectral` is less than 0.1 seconds.

## 5 Differential Chromatic Numbers of Special Graphs

**Theorem 5.** *A line graph of  $n$  nodes has differential chromatic number of  $\lfloor n/2 \rfloor$ . A cycle graph of  $n$  nodes has differential chromatic number of  $\lfloor (n-1)/2 \rfloor$ .*

*Proof.* Consider a line graph with even number of nodes labeled in order with  $n/2+1, 1, n/2+2, 2, \dots, n, n/2$ ; see Fig. 3(top). This labeling is clearly a  $n/2$ -differential coloring and it is easy to show that this is the best possible coloring. Take any labeling of this graph and consider the node labeled  $n/2$ . Regardless of the labels of its neighbor(s) this node must induce an edge difference of at most  $n/2$  achieved if the neighbor(s) is labeled  $n$ . Hence the differential chromatic number can not be more than  $n/2$ , that is,  $dc(G) = n/2 = \lfloor n/2 \rfloor$ . A similar argument can be used in the case where the graph has an odd number of vertices; see Fig. 3 (bottom).

Now consider a cycle graph with an even number of vertices labeled in order with  $1, n/2+1, 2, n/2+2, \dots, n/2, n$ ; see Fig. 4(left). This labeling is clearly  $n/2 - 1 = \lfloor (n-1)/2 \rfloor$  colorable and it is easy to show that this is the best possible coloring. Take any labeling of this graph and consider the node labeled



Table 1. Differential chromatic numbers given by `kpath` and by `GSpectral`.

name	drawing	$dc(G)$	GSpectral
petersen		3	3
hypercube		4	4
binarytree3		7	5
grid4		6	3
grid5		10	7
football		-	18



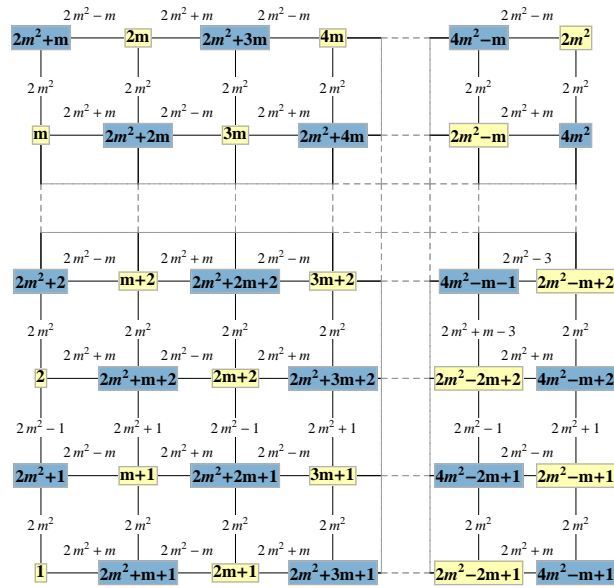


Fig. 5. Optimal coloring of  $n \times n$ -grids for even  $n$ .

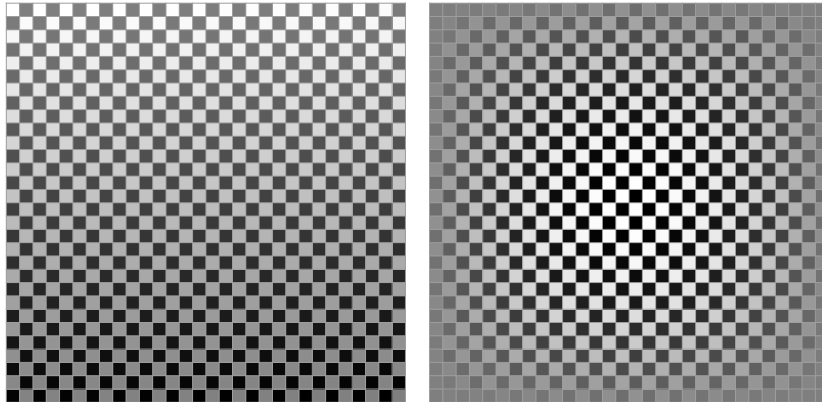


Fig. 6. Optimal grid map coloring (left) and one induced by the largest eigenvector (right).

number of isolated “dummy” nodes as there are extra colors. We further note that countries that are not neighbors, but are nevertheless close (e.g., neighbor’s neighbor), can also be forced to have distinctive colors by adding additional edges in the country graph linking these countries.

Throughout out the paper we have been assuming that the edge weights in the country graph are uniform. That is, the importance of having a different color between any pair of adjacent countries in the map is the same. We would also like to investigate ways to handle non-uniform weights, where more importance can be placed on certain countries or certain types of adjacencies (e.g., long vs. short common borders).

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